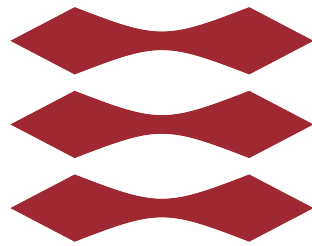


# DTU



## Discrete Mathematics - Exam Cheat Sheet

**01017**

Discrete Mathematics

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# Key Formulas & Quick Reference

## Number Theory

Formula	Description
$ab = \gcd(a, b) \cdot \text{lcm}(a, b)$	Fundamental GCD-LCM relation
$\gcd(a, b) = sa + tb$ for some $s, t \in \mathbb{Z}$	Bézout's identity
$a \equiv b \pmod{m} \iff m \mid (a - b)$	Congruence definition
$a^{-1} \pmod{m}$ exists $\iff \gcd(a, m) = 1$	Multiplicative inverse exists
$d = \gcd(a, b) \Rightarrow d^2 \mid ab$	GCD constraint on product
$a^{p-1} \equiv 1 \pmod{p}$ if $p \nmid a$	Fermat's Little Theorem
$a^p \equiv a \pmod{p}$	Fermat's Little Theorem (alt)

## Combinatorics

Formula	Description
$\binom{n}{k} = \frac{n!}{k!(n-k)!}$	Binomial coefficient
$(a + b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$	Binomial theorem
$D_n = n! \sum_{k=0}^n \frac{(-1)^k}{k!} \approx \frac{n!}{e}$	Derangements (no fixed points)
$k \cdot \binom{n}{k} = n \cdot \binom{n-1}{k-1}$	Absorption identity
$\sum_{k=0}^n \binom{n}{k} = 2^n$	Sum of binomial coefficients
$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$	Pascal's identity
$\binom{m+n}{r} = \sum_{k=0}^r \binom{m}{k} \binom{n}{r-k}$	Vandermonde's identity
Circular permutations: $(n-1)!$	Arrangements around a circle

## Graph Theory

Formula	Description
$Q_n$ : vertices = $2^n$ , edges = $n \cdot 2^{n-1}$	n-cube (hypercube)
$K_n$ : edges = $\binom{n}{2} = \frac{n(n-1)}{2}$	Complete graph
$\sum_{v \in V} \deg(v) = 2 E $	Handshaking lemma
Euler circuit exists $\iff$ all degrees even	Euler's theorem
Euler path exists $\iff$ exactly 0 or 2 odd vertices	Euler path condition
Tree on $n$ vertices has $n - 1$ edges	Tree edge count

## Set Theory & Inclusion-Exclusion

Formula	Description
$ A \cup B  =  A  +  B  -  A \cap B $	Inclusion-exclusion (2 sets)
$ A \cup B \cup C  = \sum  A_i  - \sum  A_i \cap A_j  +  A \cap B \cap C $	Inclusion-exclusion (3 sets)
$\overline{A \cap B} = \overline{A} \cup \overline{B}$	De Morgan's law
$\overline{A \cup B} = \overline{A} \cap \overline{B}$	De Morgan's law
Subsets of $n$ -element set: $2^n$	Power set cardinality
Even-sized subsets: $2^{n-1}$	Half of all subsets

## Relations

Property	Definition
Reflexive	$\forall x : (x, x) \in R$
Symmetric	$\forall x, y : (x, y) \in R \Rightarrow (y, x) \in R$
Antisymmetric	$\forall x, y : [(x, y) \in R \wedge (y, x) \in R] \Rightarrow x = y$
Transitive	$\forall x, y, z : [(x, y) \in R \wedge (y, z) \in R] \Rightarrow (x, z) \in R$
Equivalence relation	Reflexive + Symmetric + Transitive
Partial order	Reflexive + Antisymmetric + Transitive

## Examples + Solutions

### Number Theory

#### Divisibility

##### Example (Divisibility with $ab \mid cd$ )

If  $a, b, c, d$  are positive integers such that  $ab \mid cd$ , which must be true?

##### Solution:

**Key insight:**  $ab \mid cd$  does NOT imply  $a \mid c$  or  $a \mid d$  individually.

**True statement:** "If  $p$  is a prime that divides  $a$ , then  $p \mid c$  or  $p \mid d$ "

**Proof:** If  $p \mid a$  and  $ab \mid cd$ , then  $p \mid cd$ . Since  $p$  is prime,  $p \mid c$  or  $p \mid d$ .

**Counterexample:** Let  $a = 6, b = 1, c = 2, d = 3$ . Then  $ab = 6 \mid 6 = cd$ .

- But  $\gcd(a, b) = 6$  does not divide  $\gcd(c, d) = 1$
- And  $6 \nmid c$  and  $6 \nmid d$

### Example (GCD as linear combination)

Let  $a, b$  be positive integers. Which can NOT necessarily be written as  $as + bt$  for  $s, t \in \mathbb{Z}$ ?

**Solution:**

**Bézout's identity:**  $\gcd(a, b) = as + bt$  for some  $s, t \in \mathbb{Z}$ .

Any **multiple** of  $\gcd(a, b)$  can be written as  $as + bt$ .

**Answer:**  $\frac{\text{lcm}(a, b)}{\gcd(a, b)} = a \frac{b}{\gcd(a, b)^2}$  is NOT necessarily a multiple of  $\gcd(a, b)$ .

## GCD Constraints

### Example (Possible GCD values given product)

Let  $a, b$  be positive integers with  $ab = 5292 = 2^2 \cdot 3^3 \cdot 7^2$ . Which CANNOT be  $\gcd(a, b)$ ?

Options: 1, 3, 36, 42

**Solution:**

**Key fact:** If  $\gcd(a, b) = d$ , then  $d^2 \mid ab$ .

Check each:

- $d = 1$ :  $1^2 = 1 \mid 5292$  (valid)
- $d = 3$ :  $3^2 = 9 \mid 5292$  (valid, since  $3^3 \mid 5292$ )
- $d = 36 = 2^2 \cdot 3^2$ : Need  $36^2 = 2^4 \cdot 3^4 \mid 2^2 \cdot 3^3 \cdot 7^2$ . But  $2^4 \nmid 2^2$ !
- $d = 42 = 2 \cdot 3 \cdot 7$ :  $42^2 = 2^2 \cdot 3^2 \cdot 7^2 \mid 2^2 \cdot 3^3 \cdot 7^2$  (valid)

**Answer:** 36 cannot be the GCD.

## Modular Arithmetic

### Example (Congruence cancellation)

If  $ac \equiv bc \pmod{m}$ , when can we conclude  $a \equiv b \pmod{m}$ ?

**Solution:**

$ac \equiv bc \pmod{m}$  means  $m \mid c(a - b)$ .

**Cancellation Law:** If  $\gcd(c, m) = 1$ , then  $a \equiv b \pmod{m}$ .

**Counterexample when  $\gcd(c, m) \neq 1$ :**  $2 \cdot 3 \equiv 2 \cdot 6 \pmod{6}$  (both  $\equiv 0$ ), but  $3 \not\equiv 6 \pmod{6}$ .

### Example (Finding multiplicative inverses mod 9)

Find the multiplicative inverse of  $n \pmod{9}$  for  $n = 2, 6, 7$ .

**Solution:**

Inverse exists iff  $\gcd(n, 9) = 1$ .

**For  $n = 6$ :**  $\gcd(6, 9) = 3 \neq 1 \rightarrow$

Does not exist

**For  $n = 2$ :**  $\gcd(2, 9) = 1$ . Find  $x$  with  $2x \equiv 1 \pmod{9}$ :

•  $2 \cdot 5 = 10 \equiv 1 \pmod{9} \rightarrow$

**5**

**For  $n = 7$ :**  $\gcd(7, 9) = 1$ . Find  $x$  with  $7x \equiv 1 \pmod{9}$ :

•  $7 \cdot 4 = 28 \equiv 1 \pmod{9} \rightarrow$

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## Chinese Remainder Theorem

### Example (System of congruences)

Solve:  $x \equiv 1 \pmod{2}$ ,  $x \equiv 4 \pmod{5}$ ,  $x \equiv 3 \pmod{7}$

**Solution:**

Moduli 2, 5, 7 are pairwise coprime, so unique solution mod  $2 \cdot 5 \cdot 7 = 70$ .

**Method: Back substitution**

**Step 1:** From  $x \equiv 1 \pmod{2}$ :  $x = 1 + 2t_1$

**Step 2:** Substitute into  $x \equiv 4 \pmod{5}$ :  $1 + 2t_1 \equiv 4 \pmod{5} \Rightarrow 2t_1 \equiv 3 \pmod{5}$

Inverse of 2 mod 5:  $2 \cdot 3 = 6 \equiv 1$ , so  $t_1 \equiv 3 \cdot 3 = 9 \equiv 4 \pmod{5}$

Thus  $t_1 = 4 + 5t_2$ , giving  $x = 1 + 2(4 + 5t_2) = 9 + 10t_2$

**Step 3:** Substitute into  $x \equiv 3 \pmod{7}$ :  $9 + 10t_2 \equiv 3 \pmod{7} \Rightarrow 2 + 3t_2 \equiv 3 \pmod{7} \Rightarrow 3t_2 \equiv 1 \pmod{7}$

Inverse of 3 mod 7:  $3 \cdot 5 = 15 \equiv 1$ , so  $t_2 \equiv 5 \pmod{7}$

Thus  $t_2 = 5 + 7t_3$ , giving  $x = 9 + 10(5) = 59$

**Answer:**  $x \equiv 59 \pmod{70}$

**Verify:**  $59 = 29 \cdot 2 + 1$ ,  $59 = 11 \cdot 5 + 4$ ,  $59 = 8 \cdot 7 + 3$

## Functions: Injective/Surjective Analysis

### Example (Function classification)

Classify each function:

1.  $f : \mathbb{Z}^+ \rightarrow \mathbb{N}$  given by  $f(x) = \lfloor \log_2(x) \rfloor$
2.  $f : \mathbb{N} \rightarrow \mathbb{Z}$  given by  $f(x) = \begin{cases} \lfloor x/2 \rfloor & \text{if } x \text{ even} \\ -\lfloor x/2 \rfloor & \text{if } x \text{ odd} \end{cases}$
3.  $f : \mathbb{N} \rightarrow \mathbb{N}$  given by  $f(x) = x^3 + 1$

#### Solution:

1.  $f(x) = \lfloor \log_2(x) \rfloor$ ,  $\mathbb{Z}^+ \rightarrow \mathbb{N}$ :

- Surjective? Every  $n \in \mathbb{N}$  is hit by  $x = 2^n$ . Yes
- Injective?  $f(2) = f(3) = 1$ . No
- **Surjective but not injective**

2. **Alternating function**  $\mathbb{N} \rightarrow \mathbb{Z}$ :

- $f(0) = 0, f(1) = -1, f(2) = 1, f(3) = -2, f(4) = 2, \dots$
- Surjective? Hits all of  $\mathbb{Z}$ . Yes
- Injective? Each output appears exactly once. Yes
- **Bijection**

3.  $f(x) = x^3 + 1$ ,  $\mathbb{N} \rightarrow \mathbb{N}$ :

- Injective?  $x^3$  is strictly increasing. Yes
- Surjective?  $f(0) = 1, f(1) = 2, f(2) = 9, \dots$  — skips 3,4,5,6,7,8. No
- **Injective but not surjective**

## Using the Function Checker

### Example (Checking function properties with check-function)

Verify properties of  $f(x) = \lfloor \log_2(x) \rfloor$  on domain  $\{1, 2, \dots, 8\}$  onto  $\{0, 1, 2, 3\}$ :

#### Solution:

**Function f:** Injective: No, Surjective: Yes, Bijective: No

*Injectivity fails:*  $f(2) = f(3) = 1$

**Analysis:**

- Mapping:  $f(1) = 0, f(2) = 1, f(3) = 1, f(4) = 2, f(5) = 2, f(6) = 2, f(7) = 2, f(8) = 3$
- Not injective because  $f(2) = f(3) = 1$  (and others)
- Surjective because all outputs  $\{0, 1, 2, 3\}$  are hit

**More examples:**

Function	Injective	Surjective	Bijjective
$f(x) = x^2$ on $\{0, 1, 2, 3\}$ onto $\{0, 1, 4, 9\}$	Yes	Yes	Yes
$f(x) = x \bmod 5$ on $\{0, 1, 2, 3, 4\}$	Yes	Yes	Yes
$f(x) =  x $ on $\{-2, -1, 0, 1, 2\}$ onto $\{0, 1, 2\}$	No	Yes	No

## Graph Theory

### Hypercube and Complete Graphs

#### Example (Edges in $Q_n$ and $K_n$ )

**Hypercube  $Q_n$ :**

- Vertices:  $2^n$  (all  $n$ -bit binary strings)
- Each vertex has degree  $n$  (can flip any of  $n$  bits)
- By handshaking:  $2|E| = 2^n \cdot n$ , so  $|E| = n \cdot 2^{n-1}$

**Complete graph  $K_n$ :**

- Every pair of vertices connected:  $|E| = \binom{n}{2} = \frac{n(n-1)}{2}$

For  $K_{2n}$ : edges  $= \binom{2n}{2} = (2n)\frac{2n-1}{2} = n(2n-1)$

Alternative form:  $2\binom{n}{2} + n^2 = n(n-1) + n^2 = n(2n-1)$

### Degree Sequences

#### Example (Valid degree sequence?)

Does a simple graph with degrees 2, 2, 3, 3, 3, 3, 3 exist?

**Solution:**

Sum of degrees  $= 2 + 2 + 3 + 3 + 3 + 3 + 3 = 19$ .

By handshaking lemma:  $\sum \deg(v) = 2|E|$  must be even.

Since 19 is odd,

such a graph does not exist.

## Euler Circuits

### Example (Königsberg Bridge Problem)

A graph has an Euler circuit iff:

1. The graph is connected
2. Every vertex has even degree

A graph has an Euler path iff:

1. The graph is connected
2. Exactly 0 or 2 vertices have odd degree

In Königsberg: degrees are 5, 3, 3, 3 (all odd) → No Euler path or circuit.

## Combinatorics

### Binomial Theorem

#### Example (Coefficient in $(2x^2 - 3y^3)^8$ )

Find coefficients of  $x^8y^{12}$  and  $x^6y^9$ .

##### Solution:

General term:  $\binom{8}{k}(2x^2)^k(-3y^3)^{8-k} = \binom{8}{k}2^k(-3)^{8-k}x^{2k}y^{3(8-k)}$

**For  $x^8y^{12}$ :** Need  $2k = 8$  and  $3(8 - k) = 12$ .

- $k = 4$  (valid)
- Coefficient:  $\binom{8}{4} \cdot 2^4 \cdot (-3)^4 = 70 \cdot 16 \cdot 81 = 90720$

**For  $x^6y^9$ :** Need  $2k = 6$  and  $3(8 - k) = 9$ .

- $k = 3$  but  $8 - k = 5$ , and  $3 \cdot 5 = 15 \neq 9$  (invalid)
- Coefficient is 0

## Inclusion-Exclusion

### Example (Union of four sets)

Each of 4 sets has 200 elements, each pair shares 50, each triple shares 25, all four share 5. Find  $|A \cup B \cup C \cup D|$ .

**Solution:**

$$\begin{aligned}
 |A \cup B \cup C \cup D| &= \binom{4}{1} \cdot 200 - \binom{4}{2} \cdot 50 + \binom{4}{3} \cdot 25 - \binom{4}{4} \cdot 5 \\
 &= 4(200) - 6(50) + 4(25) - 1(5) = 800 - 300 + 100 - 5 = 595
 \end{aligned}$$

## Derangements

### Example (Derangement formula and values)

$$D_n = n! \sum_{k=0}^n \frac{(-1)^k}{k!} = n! \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots\right)$$

First few values:

- $D_0 = 1, D_1 = 0, D_2 = 1, D_3 = 2$
- $D_4 = 9, D_5 = 44, D_6 = 265, D_7 = 1854$

**Recurrence:**  $D_n = (n-1)(D_{n-1} + D_{n-2})$

**Approximation:**  $D_n \approx \frac{n!}{e}$  (rounds to nearest integer for  $n \geq 1$ )

## Circular Permutations

### Example (20 people around a round table)

Count seatings where two arrangements are identical if:

1. Each person has same two neighbors (ignoring direction)
2. Each person has same left AND right neighbor

**Solution:****Case 2 (same left AND right neighbor):**

- Standard circular permutation:  $(n-1)! = 19!$
- Fix one person's position, arrange remaining 19 people.

**Case 1 (same two neighbors, ignoring direction):**

- Each arrangement counted twice (clockwise vs counterclockwise)
- Answer:  $\frac{19!}{2}$

## Relations

### Example (Classify relations on $\{1, 2, 3, 4, 5, 6\}$ )

$$R_1 = \{(1, 2), (2, 3), (1, 3), (4, 5), (5, 6), (4, 6)\}$$

**Solution:**

- Reflexive? Missing  $(1, 1), (2, 2), \dots$  (no)
- Symmetric?  $(1, 2) \in R$  but  $(2, 1) \notin R$  (no)
- Antisymmetric? No pair  $(x, y), (y, x)$  with  $x \neq y$  both present (yes)
- Transitive?  $(1, 2), (2, 3) \in R$  and  $(1, 3) \in R$ ;  $(4, 5), (5, 6) \in R$  and  $(4, 6) \in R$  (yes)
- Transitive and antisymmetric only

### Example (Equivalence classes mod 4)

The equivalence relation of congruence modulo 4 on  $\mathbb{Z}$ :

$$\begin{aligned}
 [0]_{\text{mod } 4} &= \{\dots, -8, -4, 0, 4, 8, \dots\} \\
 [1]_{\text{mod } 4} &= \{\dots, -7, -3, 1, 5, 9, \dots\} \\
 [2]_{\text{mod } 4} &= \{\dots, -6, -2, 2, 6, 10, \dots\} \\
 [3]_{\text{mod } 4} &= \{\dots, -5, -1, 3, 7, 11, \dots\}
 \end{aligned} \tag{1}$$

These four equivalence classes **partition** the integers.

## Partitions of Sets

### Example (Partitions of $\mathbb{Z} \times \mathbb{Z}$ )

Which are partitions?

- (a)  $\{(x, y) : x \text{ or } y \text{ odd}\}$  and  $\{(x, y) : x \text{ and } y \text{ even}\}$   
 (b)  $\{(x, y) : x \text{ and } y \text{ odd}\}$  and  $\{(x, y) : x \text{ and } y \text{ even}\}$

#### Solution:

Every  $(x, y)$  falls into one of 4 categories: EE, OO, EO, OE

(a): " $x$  or  $y$  odd" =  $OO \cup EO \cup OE$ . " $x$  and  $y$  even" = EE.

- Disjoint? Yes. Cover everything? Yes.

• **YES, this is a partition**

(b): " $x$  and  $y$  odd" = OO. " $x$  and  $y$  even" = EE.

- Missing EO and OE!
- **NO, doesn't cover everything**

## Proof by Induction

### Example (Strong induction: Pie-throwing problem)

Prove: If  $2n + 1$  people throw pies at their nearest neighbor, at least one survives.

**Solution:**

**Base case ( $n = 1$ ):** 3 people. Closest pair (A, B) throw at each other. Third person C's nearest is either A or B. So A and B each receive one pie, C receives 0. C survives.

**Inductive step:** Assume true for  $2k + 1$  people. Consider  $2(k + 1) + 1 = 2k + 3$  people.

Let A and B be the closest pair (they throw at each other).

**Case 1:** No one else throws at A or B. The remaining  $2k + 1$  people form an independent group  $\rightarrow$  by IH, at least one survivor.

**Case 2:** At least one other person throws at A or B. Then  $\geq 3$  pies hit A or B combined. Remaining pies:  $\leq 2k + 3 - 3 = 2k$  for  $2k + 1$  people. By pigeonhole, someone survives.

**Example (Checkerboard tiling with L-triominoes)**

Every  $2^n \times 2^n$  checkerboard with one square removed can be tiled by L-triominoes.

**Solution:**

**Base case ( $n = 1$ ):**  $2 \times 2$  board with one square removed = L-triomino.

**Inductive step:** Assume true for  $2^k \times 2^k$ . For  $2^{k+1} \times 2^{k+1}$  board:

1. Divide into four  $2^k \times 2^k$  quadrants
2. The removed square is in one quadrant
3. Place one L-triomino at the center, covering one square from each of the other three quadrants
4. Now each quadrant is a  $2^k \times 2^k$  board with one square removed
5. By IH, each can be tiled.

**Polynomial Divisibility****Example ( $x^n + 1$  divisible by  $x + 1$ )**

For which positive integers  $n$  is  $x^n + 1$  divisible by  $x + 1$ ?

**Solution:**

$x + 1 \mid x^n + 1$  iff  $x = -1$  is a root of  $x^n + 1$ .

Evaluate at  $x = -1$ :  $(-1)^n + 1$

- If  $n$  odd:  $(-1)^n = -1$ , so  $-1 + 1 = 0$  (root)

- If  $n$  even:  $(-1)^n = 1$ , so  $1 + 1 = 2 \neq 0$  (not a root)

Divisible for all odd  $n$ , not divisible for any even  $n$ .

## Pigeonhole Principle

### Example (Simple graphs with all distinct degrees)

Can a simple graph on  $n \geq 2$  vertices have all distinct degrees?

#### Solution:

**Claim: NO** (by pigeonhole)

In a simple graph on  $n$  vertices:

- Possible degrees:  $0, 1, 2, \dots, n-1$  (that's  $n$  values)
- For all degrees distinct, we need exactly  $\{0, 1, 2, \dots, n-1\}$

#### But:

- Degree 0 means isolated (no neighbors)
- Degree  $n-1$  means connected to all others
- These can't coexist! (vertex with degree  $n-1$  would connect to the isolated vertex)

**Conclusion:** No simple graph on  $n \geq 2$  vertices has all distinct degrees.

## Hall's Theorem / Matching

### Example (Bipartite matching condition)

**Hall's Marriage Theorem:** A bipartite graph with parts  $X$  and  $Y$  has a matching saturating  $X$  iff for every subset  $S \subseteq X$ :

$$|N(S)| \geq |S| \quad (2)$$

where  $N(S)$  = neighbors of  $S$  in  $Y$ .

**Application:** 10 computers, 5 printers. Minimum cables so any 5 computers can print to 5 different printers?

#### Solution:

Need: every subset of 5 computers has 5 distinct printer neighbors.

If a printer connects to  $< 6$  computers, we could choose 5 computers that don't include any connected to that printer, violating Hall's condition.

Each printer must connect to  $\geq 6$  computers.

**Minimum cables:**  $5 \times 6 = 30$

## Propositional Logic (Truth Sayer/Liar Puzzles)

### Example (Truth Sayer and Liar Logic)

Peter says: "At least one of us is a liar." What are Peter and Signe?

#### Solution:

Let  $P$  = "Peter is truth sayer",  $S$  = "Signe is truth sayer"

Peter's claim:  $\neg P \vee \neg S$

**Key:** If Peter is a truth sayer, his claim must be true. If he's a liar, his claim must be false.

$$P \Leftrightarrow (\neg P \vee \neg S) \quad (3)$$

$P$	$S$	$\neg P \vee \neg S$	$P \Leftrightarrow (\neg P \vee \neg S)$
T	T	F	F
T	F	T	<b>T</b>
F	T	T	F
F	F	T	F

**Answer:** Peter is a truth sayer, Signe is a liar.

## Perfect Numbers

### Example (Verify perfect numbers)

Show that 6 and 28 are perfect numbers (equal to sum of proper divisors).

#### Solution:

**For 6:** Divisors (excluding 6): 1, 2, 3

$$1 + 2 + 3 = 6 \quad (4)$$

**For 28:** Divisors (excluding 28): 1, 2, 4, 7, 14

$$1 + 2 + 4 + 7 + 14 = 28 \quad (5)$$

**Theorem:**  $2^{p-1}(2^p - 1)$  is perfect when  $2^p - 1$  is prime (Mersenne prime).

Example:  $p = 3$ ,  $2^3 - 1 = 7$  (prime), so  $2^2 \cdot 7 = 28$  is perfect.

## Set Operations Proofs

**Example (Prove  $(A - C) \cap (C - B) = \emptyset$ )**

**Solution:**

$(A - C)$  = elements in  $A$  but not in  $C$

$(C - B)$  = elements in  $C$  but not in  $B$

For  $x \in (A - C) \cap (C - B)$ :

- $x \in A - C$  means  $x \in A$  and  $x \notin C$
- $x \in C - B$  means  $x \in C$  and  $x \notin B$

**Contradiction:**  $x \notin C$  and  $x \in C$  cannot both be true.

Therefore  $(A - C) \cap (C - B) = \emptyset$ .

**Example (Prove  $(B - A) \cup (C - A) = (B \cup C) - A$ )**

**Solution:**

**LHS:**  $x \in (B - A) \cup (C - A)$

- $x \in B - A$  or  $x \in C - A$
- $(x \in B \wedge x \notin A)$  or  $(x \in C \wedge x \notin A)$
- $(x \in B \vee x \in C)$  and  $x \notin A$

**RHS:**  $x \in (B \cup C) - A$

- $x \in B \cup C$  and  $x \notin A$
- $(x \in B \vee x \in C)$  and  $x \notin A$

Both sides are equivalent.

## Equivalence Relations

**Example (Cardinality as equivalence relation)**

Let  $R$  on sets of real numbers:  $SRT$  iff  $|S| = |T|$ .

**Solution:**

**Reflexive:**  $|S| = |S|$  (yes)

**Symmetric:**  $|S| = |T| \Rightarrow |T| = |S|$  (yes)

**Transitive:**  $|S| = |T|$  and  $|T| = |U| \Rightarrow |S| = |U|$  (yes)

This is an equivalence relation.

**Equivalence classes:**

- $\{\{0, 1, 2\}\}$  = all sets with exactly 3 elements
- $[\mathbb{Z}]$  = all countably infinite sets (includes  $\mathbb{N}$ ,  $\mathbb{Q}$ )

**Example (Rational equivalence:**  $(a, b)R(c, d)$  **iff**  $ad = bc$ )

**Solution:**

This is an equivalence relation (represents fractions  $\frac{a}{b} = \frac{c}{d}$ ).

**Reflexive:**  $a \cdot b = b \cdot a$  (yes)

**Symmetric:**  $ad = bc \Rightarrow cb = da$  (yes)

**Transitive:** If  $ad = bc$  and  $cf = de$ , then:

- Multiply:  $adf = bcf = bde$
- Since  $d > 0$ :  $af = be$  (yes)

This is an equivalence relation.

Equivalence class of  $(1, 2)$ : all pairs  $(k, 2k)$  for  $k \in \mathbb{Z}^+$

## Generalized Pigeonhole

**Example (Generalized pigeonhole for  $n$  boxes)**

If  $n_1 + n_2 + \dots + n_t - t + 1$  objects are placed in  $t$  boxes, then some box  $i$  contains at least  $n_i$  objects.

**Solution:**

**Proof by contradiction:**

Assume each box  $i$  contains fewer than  $n_i$  objects (at most  $n_i - 1$ ).

Total objects  $\leq (n_1 - 1) + (n_2 - 1) + \dots + (n_t - 1) = \sum n_i - t$

But we have  $\sum n_i - t + 1$  objects.

$\sum n_i - t + 1 \leq \sum n_i - t$  implies  $1 \leq 0$ . Contradiction!

## Polynomial Division (using auto-div)

### Example (Polynomial division with auto-div)

Divide  $x^4 + 3x^3 + \frac{5}{2}x + 6$  by  $x + 2$ :

$$\begin{array}{r}
 \phantom{x+2} \overline{x^3 \phantom{+} x^2 \phantom{-} 2x + \frac{13}{2}} \\
 x+2 \overline{x^4 + 3x^3 \phantom{+} \frac{5}{2}x + 6} \\
 \underline{-(x^4 + 2x^3)} \phantom{+} \\
 \phantom{x+2} x^3 \phantom{+} \frac{5}{2}x + 6 \\
 \underline{-(x^3 + 2x^2)} \phantom{+} \\
 \phantom{x+2} -2x^2 + \frac{5}{2}x + 6 \\
 \underline{-(-2x^2 - 4x)} \phantom{+} \\
 \phantom{x+2} \phantom{-} \frac{13}{2}x + 6 \\
 \underline{-(\frac{13}{2}x + 13)} \\
 \phantom{x+2} \phantom{-} \phantom{\frac{13}{2}x} -7
 \end{array} \tag{6}$$

## Calculation Workspace

### Quick Reference: Built-in Typst Functions

Function	Example	Result
calc.gcd(a, b)	calc.gcd(48, 18)	6
calc.lcm(a, b)	calc.lcm(12, 18)	36
calc.fact(n)	calc.fact(6)	720
calc.binom(n, k)	calc.binom(10, 3)	120
calc.perm(n, k)	calc.perm(5, 3)	60
calc.rem(a, b)	calc.rem(17, 5)	2
calc.quo(a, b)	calc.quo(17, 5)	3
calc.pow(a, b)	calc.pow(2, 10)	1024

### Binomial Coefficients

$$\binom{10}{5} = 252 \tag{7}$$

$$\binom{8}{4} = 70 \tag{8}$$

$$\binom{20}{10} = 184756 \quad (9)$$

$$\binom{15}{7} = 6435 \quad (10)$$

## Factorials

$$5! = 120 \quad (11)$$

$$7! = 5040 \quad (12)$$

$$10! = 3628800 \quad (13)$$

## Derangements

$$D_4 = 9 \quad (14)$$

$$D_5 = 44 \quad (15)$$

$$D_6 = 265 \quad (16)$$

$$D_7 = 1854 \quad (17)$$

## GCD and LCM

$$\gcd(48, 18) = 6 \quad (18)$$

$$\gcd(5292, 36) = 36 \quad (19)$$

$$\gcd(662, 414) = 2 \quad (20)$$

$$\text{lcm}(12, 18) = 36 \quad (21)$$

$$\text{lcm}(24, 36) = 72 \quad (22)$$

## Bézout Coefficients (Extended Euclidean Algorithm)

$$\gcd(48, 18) = 6 = (48)(-1) + (18)(3) \quad (23)$$

$$\gcd(35, 15) = 5 = (35)(1) + (15)(-2) \quad (24)$$

$$\gcd(662, 414) = 2 = (662)(-5) + (414)(8) \quad (25)$$

## Modular Inverses

$$2^{-1} \equiv 5 \pmod{9} \quad (26)$$

$$6^{-1} \pmod{9} \quad (27)$$

does not exist (since  $\gcd(6, 9) = 3 \neq 1$ )

$$7^{-1} \equiv 4 \pmod{9} \quad (28)$$

$$3^{-1} \equiv 5 \pmod{7} \quad (29)$$

$$5^{-1} \equiv 5 \pmod{12} \quad (30)$$

## Chinese Remainder Theorem

**CRT Solution:**  $x \equiv 59 \pmod{70}$

Verify:  $59 \equiv 1 \pmod{2}$ ,  $59 \equiv 4 \pmod{5}$ ,  $59 \equiv 3 \pmod{7}$

**CRT Solution:**  $x \equiv 68 \pmod{105}$

Verify:  $68 \equiv 2 \pmod{3}$ ,  $68 \equiv 3 \pmod{5}$ ,  $68 \equiv 5 \pmod{7}$

## Graph Theory Quick Calculations

$$Q_4 : 16 \text{ vertices, } 32 \text{ edges} \quad (31)$$

$$Q_5 : 32 \text{ vertices, } 80 \text{ edges} \quad (32)$$

$$K_6 : 6 \text{ vertices, } 15 \text{ edges} \quad (33)$$

$$K_{10} : 10 \text{ vertices, } 45 \text{ edges} \quad (34)$$

$$K_{20} : 20 \text{ vertices, } 190 \text{ edges} \quad (35)$$

## Inclusion-Exclusion (4 sets with equal intersections)

$$|A \cup B \cup C \cup D| = 4(200) - 6(50) + 4(25) - 5 = 595 \quad (36)$$

## Primality and Primes

$$17 \text{ is prime} \quad (37)$$

$$91 \text{ is not prime} \quad (38)$$

$$97 \text{ is prime} \quad (39)$$

There are 10 primes below 30: 2, 3, 5, 7, 11, 13, 17, 19, 23, 29

## General Linear Congruences

Solve  $ax \equiv c \pmod{m}$ :

$$3x \equiv 6 \pmod{9} \Rightarrow x \equiv 2 \pmod{3} \quad (40)$$

$$4x \equiv 5 \pmod{9} \Rightarrow x \equiv 8 \pmod{9} \quad (41)$$

$$6x \equiv 15 \pmod{21} \Rightarrow x \equiv 6 \pmod{7} \quad (42)$$

## Division with Remainder

$$17 = 5 \cdot 3 + 2 \quad (43)$$

$$100 = 7 \cdot 14 + 2 \quad (44)$$

$$5292 = 36 \cdot 147 + 0 \quad (45)$$

## Relation Properties

**Relation  $R_1$  on set (1, 2, 3):**

- Reflexive: Yes
- Symmetric: Yes
- Antisymmetric: No
- Transitive: Yes
- **Equivalence relation**

**Relation  $R_2$  on set (1, 2, 3):**

- Reflexive: Yes
- Symmetric: No
- Antisymmetric: Yes
- Transitive: Yes
- **Partial order**

## Function Property Checker

Check if functions are injective/surjective/bijective on finite domains:

Usage	Code
Define function	<pre>#let my_func = (x) =&gt; calc.floor(calc.log(x, base: 2))</pre>
Check properties	<pre>#let result = check-function(   my_func,   (1, 2, 3, 4, 5, 6, 7, 8), // domain   codomain: (0, 1, 2, 3)   // codomain (optional) )</pre>
Display results	<pre>#show-function-check(result, func-name: "f")</pre>

**Quick examples:**

- $f(x) = \lfloor \log_2(x) \rfloor$ : Inj: No, Surj: Yes, Bij: No
- $g(x) = x^2$  on  $\{0, 1, 2, 3\}$ : Inj: Yes, Surj: Yes, Bij: Yes
- $h(x) = |x|$  on  $\{-2, \dots, 2\}$ : Inj: No, Surj: Yes, Bij: No

**Note:** Only works for finite domains. For infinite domains ( $\mathbb{Z}$ ,  $\mathbb{N}$ ,  $\mathbb{R}$ ), use mathematical proofs.

## Your Calculations Here